Applications of Graph Theory to Covering Arrays

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joint work with Lucia Moura and Brett Stevens

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Covering Arrays

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tests	bedroom	hall	bathroom	kitchen
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4	1	1	0	1
5	1	1	1	0

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Any two columns are *qualitatively independent* (every possible pair occurs in some row.)

A covering array on a graph G, denoted CA(n, G, k), is:

- \star an $n \times r$ array where r = |V(G)|,
- \star with entries from \mathbb{Z}_k (k is the alphabet),
- * columns for adjacent vertices are qualitatively independent (have all possible pairs).

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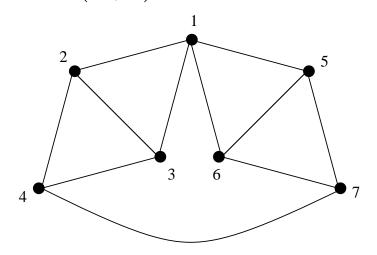
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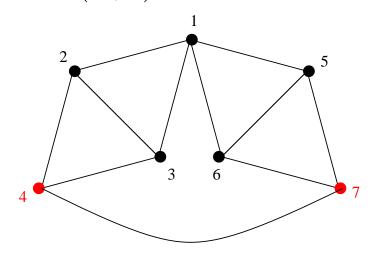


1	2	3	4	5	6	7
0	0	0	0	0	0	0
0	1	1	0	1	1	1
1	0	1	1	0	1	1
1	1	0	1	1	0	1
1	1	1	1	1	1	0

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Graph Homomorphisms

A homomorphism of graphs G and H is a map

$$f:V(G)\to V(H)$$

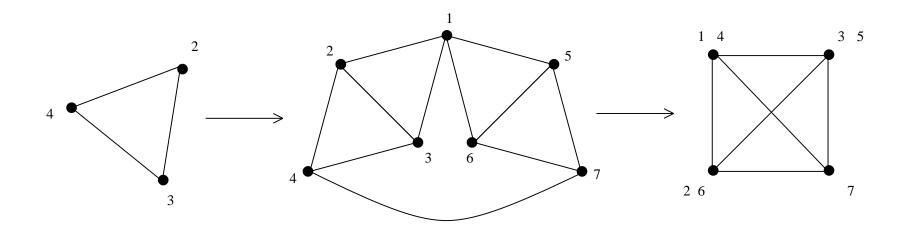
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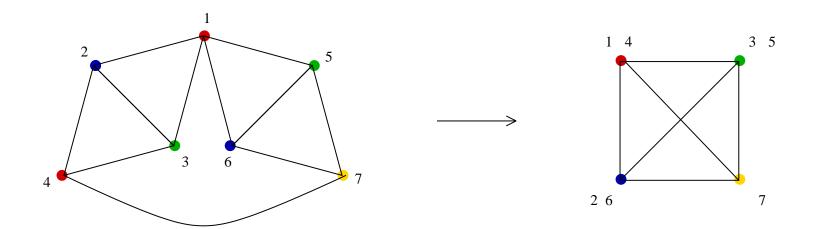
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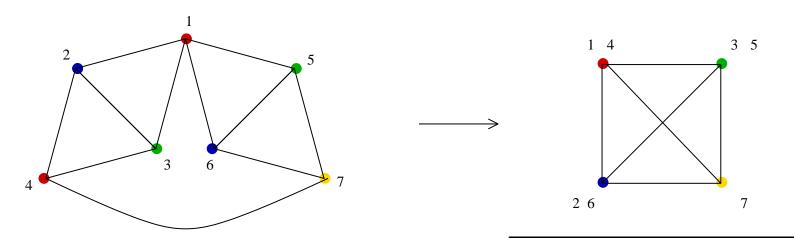
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A Construction

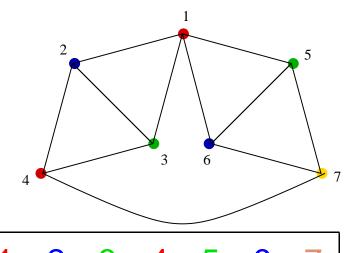


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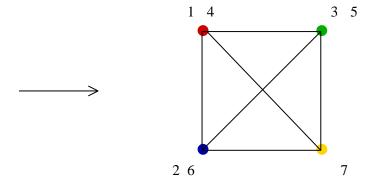


red	red blue green yellow					
0	0	0	0			
0	1	1	1			
1	0	1	1			
1	1	0	1			
1	1	1	0			

A Construction



1	2	3	4	5	6	7
0	0	0	0	0	0	0
0	1	1	0	1	1	1
1	0	1	1	0	1	1
1	1	0	1	1	0	1
1	1	1	1	1	1	0



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There is a lower bound from maximum clique size

$$CAN(K_{\omega(G)}, k) \leq CAN(G, k) \leq CAN(K_{\chi(G)}, k),$$

and an upper a bound from the chromatic number.

Partitions

	1	2	3	4
1	0	0	0	0
2	0	1	1	1
1 2 3 4	0	2	2	
4	1	0	1	2 2 0
5	1	1	1 2	0
6 7	1	2	0	1
7	2	0	2	1
8	1 1 2 2 2	1	0	2
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$$1 \to \{\{1, 2, 3\}, \{4, 5, 6\}, \{7, 8, 9\}\}\}$$

$$2 \to \{\{1, 4, 7\}, \{2, 5, 8\}, \{3, 6, 9\}\}\}$$

$$3 \to \{\{1, 6, 8\}, \{2, 4, 9\}, \{3, 5, 7\}\}\}$$

$$4 \to \{\{1, 5, 9\}, \{2, 6, 7\}, \{3, 4, 8\}\}$$

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8	1 2 2 2	1	0	2
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The columns of a covering array with a k-alphabet and n rows determine k-partitions of an n-set.

Qualitative Independence

Let A, B be k-partitions of an n-set,

$$A = \{A_1, A_2, \dots, A_k\}$$
 and $B = \{B_1, B_2, \dots, B_k\}$.

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Two partitions formed from the columns in a covering array are qualitatively independent.

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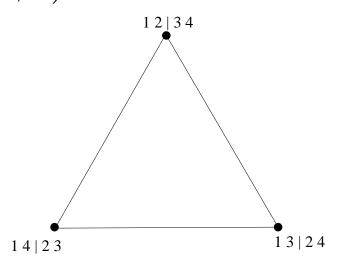
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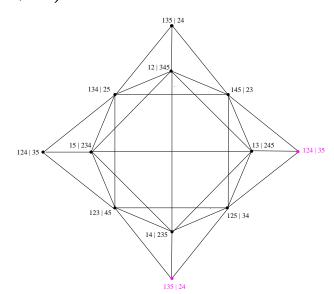
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The graph QI(5,2):



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Theorem (M. and Stevens, 2002) The minimal size of a covering array on a graph G is

$$\min\{n:G\to QI(n,k).\}$$

The uniform qualitative independence graph, $UQI(k\ell,k)$ has

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A clique in $UQI(k\ell,k)$ is equivalent to a balanced covering array (each letter occurs the same number of times in each column.)

Example 1: $UQI(k^2, k)$

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- \star This bounds the size of a clique by k+1.
- * A covering array with k^2 rows on a k-alphabet can have no more than k+1 columns.

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UQI(12,3) in an association scheme so the ratio bound holds

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There does not exist a balanced covering array with 12 rows and 8 columns with an alphabet of size three.

Open problems:

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- \star Find all eigenvalues of $UQI(k\ell,k)$.
- \star Generalize eigenvalue bound for UQI(12,3) to more cases.